

Generalizations of Thompson Group V using Synchronizing Automata

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n -ary Cantor Space

$\mathcal{C}_n = \{0, 1, \dots, n-1\}^\omega$, the infinite words over an n letter alphabet

Let v_1, v_2, \dots, v_k be a list of finite prefixes that 'partition' \mathcal{C}_n , i.e. a maximal antichain with respect to lexicographical order.

Also let u_1, u_2, \dots, u_k be another maximal antichain of the same length.

Build the function $f : \mathcal{C}_n \rightarrow \mathcal{C}_n$ where $f(v_i w) = u_i w$ for each i . We call f a finite prefix exchange.

Example: 0, 10, 11 and 01, 1, 00

Remark

Each finite prefix exchange has infinitely many representations.

Generalized Thompson Groups V_n

The group V_n is the group of all finite prefix exchanges on \mathcal{C}_n .

- ▶ V_2 is commonly known as V .
- ▶ V_n is finitely presentable.
- ▶ V_n is simple when n is even.
- ▶ V_n has simple derived subgroup of index two when n is odd.
- ▶ $V_n \cong V_m$ iff $n = m$.

Subgroup Generating Relations

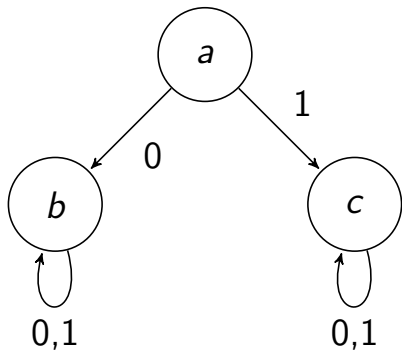
Let \sim be an equivalence relation on $\{0, 1, \dots, n-1\}^*$, the set of finite prefixes. We call \sim subgroup generating when for all finite prefixes α and β ,

$\alpha \sim \beta$ if and only if

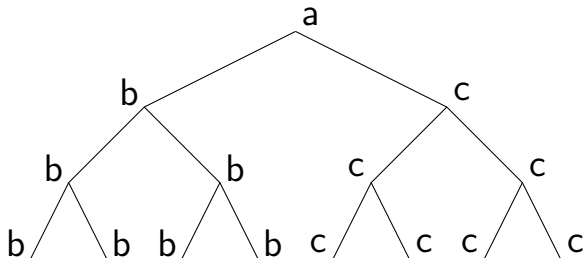
for all $i \in \{0, 1, \dots, n-1\}$, $\alpha i \sim \beta i$.

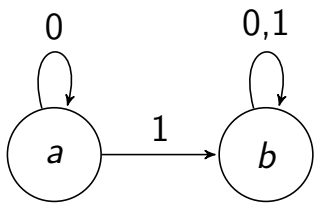
Relation Preserving Subgroups of V_n

We denote the subgroup of V_n that preserves \sim by V_{\sim} .

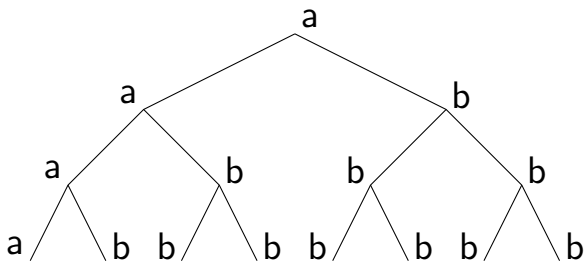


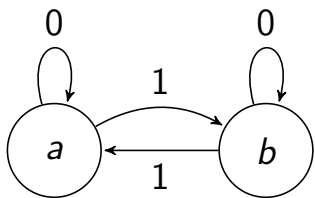
$$V_{\sim} \cong V \times V$$



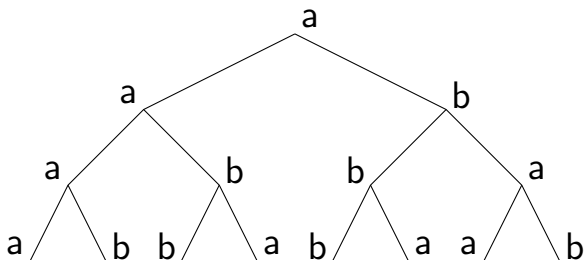


$$V_{\sim} \cong \text{Stab}_V(\bar{0})$$





$$V_{\sim} \cong V$$



Theorem (Bleak, Cameron, Maissel, Navas, Olukoya)

Automorphisms of V_n are precisely conjugations of V_n by homeomorphisms of C_n induced by bi-synchronizing automata.

Synchronizing Automata

An automaton A is synchronizing (at level k) if every word of length k is a reset word. Furthermore, A is bi-synchronizing if A and A^{-1} are both synchronizing.

Lemma

$V_n^A \leq V_n$ if and only if A is synchronizing.

Theorem (D, Olukoya)

Let A be synchronizing. Then $V_n^A = V_{\sim}$ where \sim is constructed using A^{-1} .