

SandGAL 2019

"Semigroups and Groups, Automata, Logics" | 10-13 June, 2019

A REDUCTION THEOREM FOR NONSOLVABLE FINITE GROUPS

(joint w. F. Leinen & O. Puglisi).

Let G be a finite group.

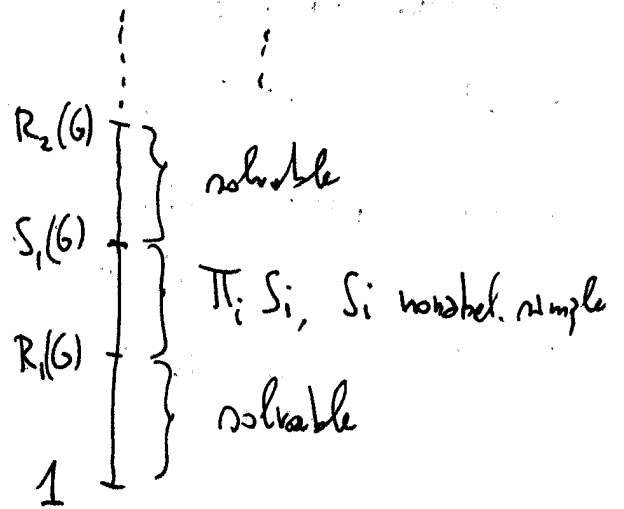
$R(G) := \langle H \mid H \text{ solvable normal subgroup of } G \rangle$

$S(G) := \langle K \mid K \text{ minimal normal nonabelian subgroup of } G \rangle$

The "R-S-Series" of G :

$$\begin{cases} R_1(G) = R(G) \\ \frac{S_i(G)}{R_i(G)} = S\left(\frac{G}{R_i(G)}\right), \forall i \geq 1 \\ \frac{R_{i+1}(G)}{S_i(G)} = R\left(\frac{G}{S_i(G)}\right), \forall i \geq 1 \end{cases}$$

$$R_{n+1}(G) = G$$



If $G = R_{n+1}(G) > R_n(G)$, then $n :=$ the nonsolvable length of $G = \lambda(G)$

$\lambda(G) = 0 \iff G$ is solvable.

$\lambda(G) = 1 \iff G$ (direct product of) simple / almost simple / quasimple qps.

$\lambda(\underbrace{(\dots (A_5 \wr A_5) \wr A_5) \wr \dots \wr A_5}_{n\text{-times}}) = n$.

- E. Khukhro & P. Shumyatsky (2015)

$(\lambda(G), \lambda_p(G) \text{ also}).$

Theorem (F.L.P. 2019)

Every finite group of nonsolvable length n contains an n -RAREFIED subgroup.

n -RAREFIED GROUPS:

G n.t. $\lambda(G) = n$ & the following holds:

1) $R_1(G) = \phi(G) \cong \frac{R_{i+1}(G)}{S_i(G)} = \phi\left(\frac{G}{S_i(G)}\right) \quad \forall i \geq 1;$

2) Each $\frac{S_i(G)}{R_i(G)}$ is the unique minimal normal subgp. in $\frac{G}{R_i(G)}$, $\forall i \geq 1$.

3) The components of each $\frac{S_i(G)}{R_i(G)}$ lie in

$$\mathcal{L} = \left\{ \begin{array}{l} \text{PSL}_2(q), \text{ for } q \in \{2^r, 3^r, p^{2^a}\}, \text{ with } r, p \in \mathbb{P}, 2 \neq p, a \geq 0; \\ \text{PSL}_3(3), {}^2B_2(2^a), \text{ } a \in \mathbb{P} \end{array} \right\}.$$

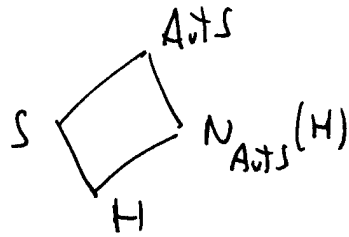
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Note that $\mathcal{L} \supset \mathcal{M} = \{\text{minimal simple gps}\}$, ($A_G \in \mathcal{L} \cup \mathcal{M}$).

Lemma. Assume that S is a finite simple (nonabelian) group.
If $S \notin \mathcal{L}$ then $\exists H \leq S$ n.t.

- 1) $\lambda(H) = 1$
- 2) $H = N_S(H)$
- 3) $\text{Aut}(S) = S \cdot N_{\text{Aut}(S)}(H)$



Applications of the Thm

$\lambda(G) \leq \max \{ \lambda_2(A) \mid A \text{ soluble subgp of } G \}$.

If $\lambda(G) \geq 1$, then $\exists H \leq G$, H 2-generated n.t. $\lambda(H) = \lambda(G)$.

$\max \{ \lambda(G) \mid G \leq \text{Sym}(m) \} = \lfloor \log_5(m) \rfloor$

$\max \{ \lambda(G) \mid G \leq \text{GL}_m(\mathbb{F}) \} = 1 + \lfloor \log_5(m/2) \rfloor$ (if $|\mathbb{F}| \geq 4$
when $|\mathbb{F}| \leq 3$ bounds)

Conjecture I Let G n.t. $\lambda(G) = n$ and w a ^{reduced} word in $k \geq 2$ variables.
If $\text{length}(w) = n$, then $w(G) \neq 1$.

Conjecture II
("The Puzlisi Alternative")

Any fin. gen. profinite group either contains a free subgroup or admits a prosolvable group which is "big enough".