

The intersection problem in $F_n \times \mathbb{Z}^m$

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Semigroups and Groups, Automata, Logics

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(joint with J. Delgado)

Outline

- 1 Free-abelian-by-free groups
- 2 Stallings' automata
- 3 Vectored Stallings' automata
- 4 Membership
- 5 Intersection

Finitely generated free-abelian groups

This talk is about semidirect products of f.g. free-abelian groups by f.g. free groups ([free-abelian-by-free](#), for short).

For all the talk,

- *elements from \mathbb{Z}^m are [row](#) integral vectors of length m , with [additive](#) notation;*
- *think [endomorphisms](#) of \mathbb{Z}^m as $m \times m$ integral matrices \mathbf{A} , acting on the [right](#) $\mathbf{u} \mapsto \mathbf{uA}$;*
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Abelian-by-free groups

- Let $X = \{x_1, \dots, x_n\}$ be a set, and \mathbb{F}_X the free group on it;
- fix matrices $\mathbf{A}_j \in \text{GL}_m(\mathbb{Z})$, $j = 1, \dots, n$ and consider the action $\mathbf{A}_\bullet : \mathbb{F}_X \rightarrow \text{Aut}(\mathbb{Z}^m)$, $x_j \mapsto \mathbf{A}_j$;
- consider the corresponding semidirect product $\mathbb{G}_\mathbf{A} = \mathbb{F}_n \rtimes_{\mathbf{A}_\bullet} \mathbb{Z}^m$,

$$\mathbb{G}_\mathbf{A} = \left\langle \begin{array}{c|c} x_1, \dots, x_n & \begin{array}{l} t_i t_k = t_k t_i \quad \forall i, k \in [1, m] \\ x_j^{-1} t_i x_j = t_i^{\mathbf{e}_j \mathbf{A}_j} \quad \forall i \in [1, m], \forall j \in [1, n] \end{array} \end{array} \right\rangle$$

- Notation: $t^{\mathbf{u}} := t^{(u_1, \dots, u_m)} := t_1^{u_1} t_2^{u_2} \dots t_m^{u_m}$, where 't' is meaningless;
- This way, we get multiplicative notation, $t^{\mathbf{u}} t^{\mathbf{v}} = t^{\mathbf{u} + \mathbf{v}}$.

Of course, the case of *trivial* action, $\mathbf{A}_j = \mathbf{I}_m$, corresponds to the direct product $\mathbb{F}_X \times \mathbb{Z}^m$, where the x_j 's commute with the t_i 's.

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Observation

We have the standard split short exact sequence,

$$1 \longrightarrow \mathbb{Z}^m \longrightarrow \mathbb{G}_{\mathbf{A}} = \mathbb{F}_X \rtimes_{\mathbf{A}} \mathbb{Z}^m \xrightarrow{\pi} \mathbb{F}_X \longrightarrow 1.$$

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Writing the semidirect relation as $t^u x_j = x_j t^{uA_j}$ or $t^u x_j^{-1} = x_j^{-1} t^{uA_j^{-1}}$, we get normal forms for elements $g \in \mathbb{G}_{\mathbf{A}}$ as

$$g = wt^{\mathbf{u}} = w \cdot t_1^{u_1} t_2^{u_2} \cdots t_m^{u_m},$$

where $w = g\pi \in \mathbb{F}_X$ and $\mathbf{u} = (u_1, u_2, \dots, u_m) \in \mathbb{Z}^m$.

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Subgroups of free-abelian-by-free groups

Proposition

Let $\mathbb{G}_{\mathbf{A}} = \mathbb{F}_X \rtimes_{\mathbf{A}_\bullet} \mathbb{Z}^m$. Then, any subgroup $H \leq \mathbb{G}_{\mathbf{A}}$ admits the decomposition $H \simeq H\pi \rtimes L_H$, where $H\pi \leq \mathbb{F}_X$, $L_H = H \cap \mathbb{Z}^m \leq \mathbb{Z}^m$, and the action is the restriction of \mathbf{A}_\bullet to both $H\pi$ and L_H , i.e., $H\pi \rightarrow \text{Aut}(L_H)$, $w \mapsto \mathbf{A}_w|_{L_H}$.

Corollary

Any subgroup H of a free-abelian-by-free group $\mathbb{G}_{\mathbf{A}}$ is again free-abelian-by-free. Moreover, H is finitely generated if and only if $H\pi$ is finitely generated.

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To *extend* the Stallings bijection

$$\begin{array}{ccc}
 \{\text{Subgroups of } \mathbb{F}_X\} & \longleftrightarrow & \{\text{Stallings } X\text{-automata}\} \\
 H & \mapsto & \Gamma(H) \\
 L(\mathcal{A}) & \leftarrow & \mathcal{A}
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to *subgroups of* \mathbb{G}_A , and use it to *solve several algorithmic problems* in free-abelian-by-free groups. We shall focus on:

- *the membership problem;* ✓
- *the subgroup conjugacy problem;* *caution! there are $F_{14} \rtimes \mathbb{Z}^4$ groups with unsolvable CP*
- *the intersection problem;* *caution! $F_2 \times \mathbb{Z}$ is NOT Howson*

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- 2 **Stallings' automata**
- 3 Vectored Stallings' automata
- 4 Membership
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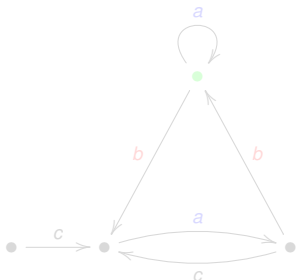
Stallings' automata

Definition

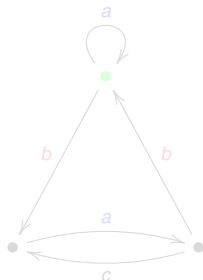
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- 1- it is *connected*,
- 2- it is *trim*, (no vertex of degree 1 except possibly q_0),
- 3- it is *deterministic* (no two edges with the same label go out of (or into) the same vertex).

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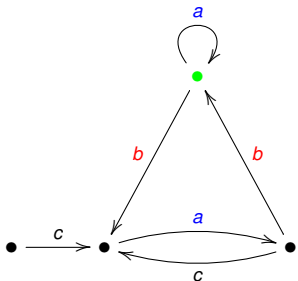
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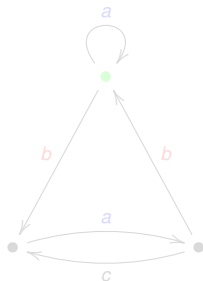
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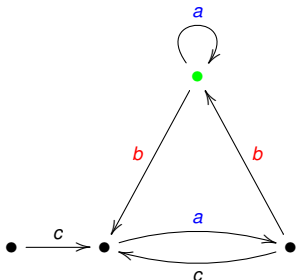
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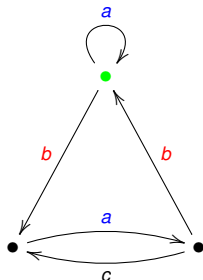
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Stallings (building on previous works) gave a **bijection** between finitely generated subgroups of F_X and Stallings automata:

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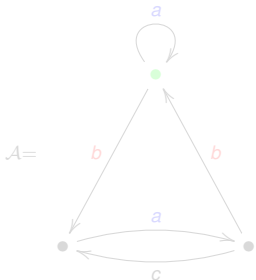
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Reading the subgroup from the automata

Definition

To any given Stallings automaton $\mathcal{A} = (V, E, q_0)$, we associate its language:

$$L(\mathcal{A}) = \{ \text{labels of closed paths at } q_0 \} \leq F(A).$$



$$L(\mathcal{A}) = \{1, a, a^{-1}, bab, bc^{-1}b, babab^{-1}cb^{-1}, \dots\}$$

$$L(\mathcal{A}) \not\ni bc^{-1}bcaa$$

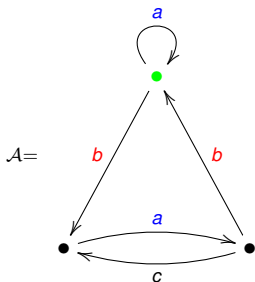
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A basis for $L(\mathcal{A})$

Proposition

For every Stallings automaton $\mathcal{A} = (V, E, q_0)$, and every maximal tree T , the group $L(\mathcal{A})$ is free with free basis

$$\{x_e = \text{label}(T[q_0, \iota e]) \cdot e \cdot T[\tau e, q_0] \in L(\mathcal{A}) \mid e \in EX - ET\},$$

where $T[p, q]$ denotes the geodesic in T from p to q . In particular, $\text{rk}(L(\mathcal{A})) = 1 - |V| + |E|$.

Constructing the automaton from the subgroup

Given $H = \langle w_1, \dots, w_n \rangle \in F(A)$, construct the *flower automaton*, denoted $\mathcal{F}(H)$.

Clearly, $L(\mathcal{F}(H)) = H$.

... But $\mathcal{F}(H)$ is not in general deterministic...

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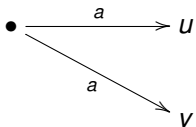
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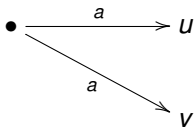
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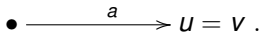
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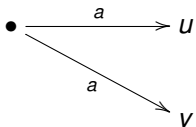
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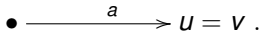
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Constructing the automata from the subgroup

Lemma (Stallings)

If $\mathcal{A} \rightsquigarrow \mathcal{A}'$ is a Stallings folding then $L(\mathcal{A}) = L(\mathcal{A}')$.

Given a f.g. subgroup $H = \langle w_1, \dots, w_n \rangle \leq F_X$ (we assume w_i are reduced words), do the following:

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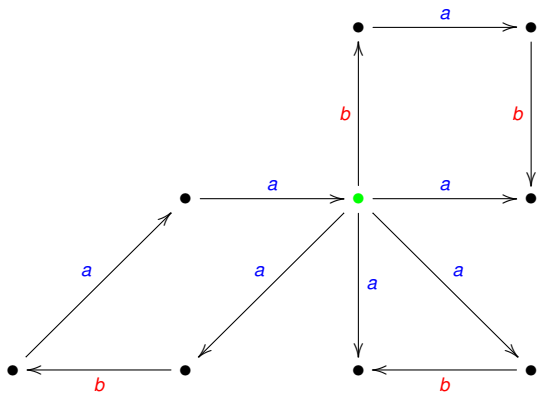
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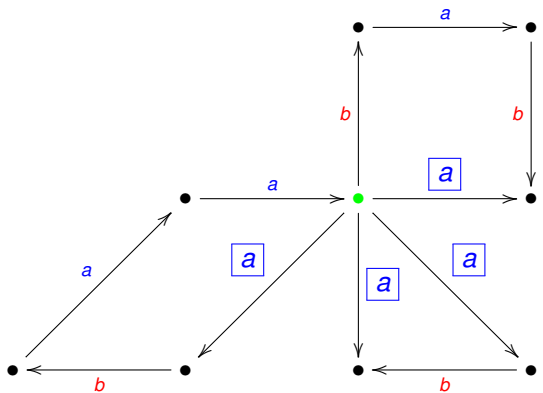
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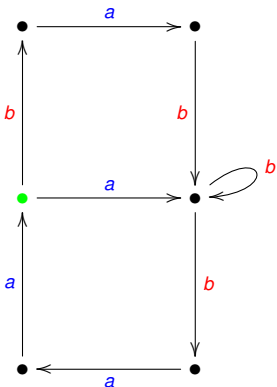
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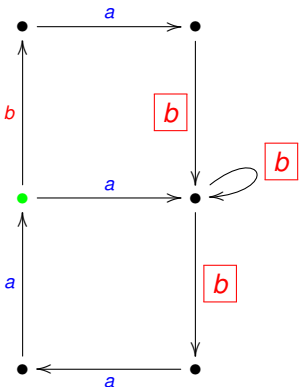
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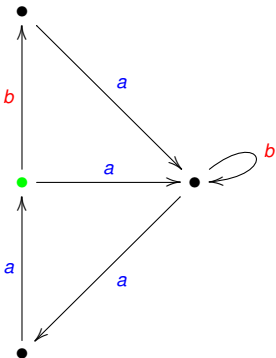
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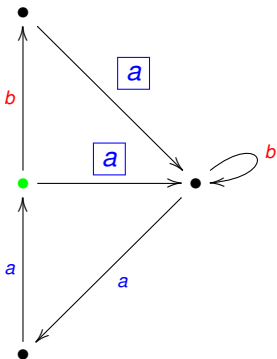
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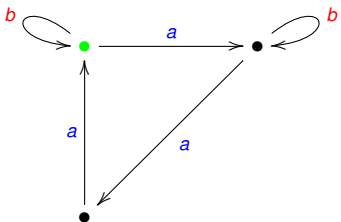
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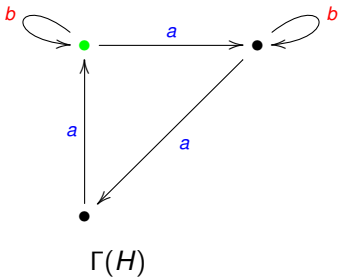


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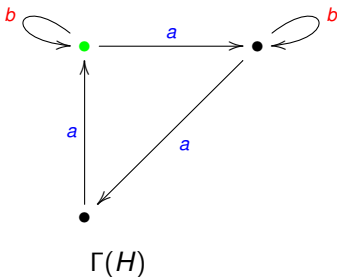
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Local confluence

It can be shown that

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The automaton $\Gamma(H)$ does not depend on the sequence of foldings.

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The automaton $\Gamma(H)$ does not depend on the generators of H .

Theorem

The following is a well defined bijection:

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- 3 Vectored Stallings' automata**
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Vectored Stallings' automata

Definition

Let us consider now, *vectored X-automata*, i.e., X-graphs with *vectors* assigned at the heads and tails of the edges,

$$\bullet \xrightarrow{u_1 \quad a \quad u_2} \bullet ,$$

reading $t^{-u_1} a t^{u_2} = a t^{u_2 - u_1} A$ (and the inverse if traversed backwards).
... plus a *subspace* $L \leq \mathbb{Z}^m$ attached to the basepoint.

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For a f.g. subgroup $H = \langle w_1 t^{u_1}, \dots, w_r t^{u_r} \rangle$ of $G_A = \mathbb{F}_n \rtimes_{A_1, \dots, A_n} \mathbb{Z}^m$, we can also construct a *flower automaton*.

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Abelian moves

Definition

We need now some extra operations to allow moving abelian mass around:

- *edge moves,*
- *vertex moves,*
- *vertex moves at the basepoint,*
- *open foldings,*
- *closed foldings,*
- *increase L to its closure by the labels of all closed paths at \bullet .*

Definition

A *vectored Stallings A -automata* is a connected and trim vectored A -automata satisfying:

- \mathcal{A} is deterministic,
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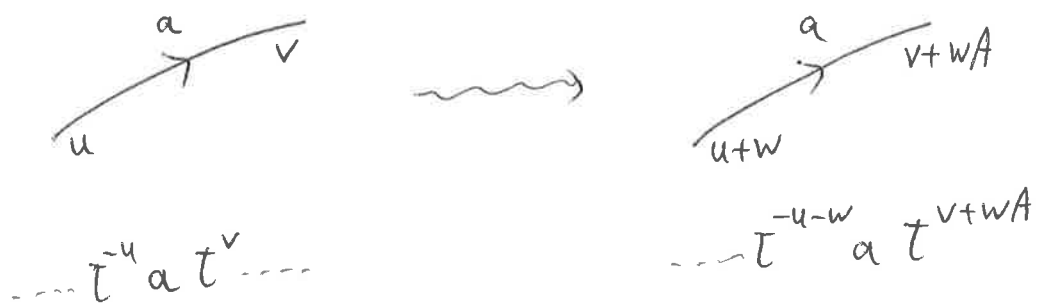
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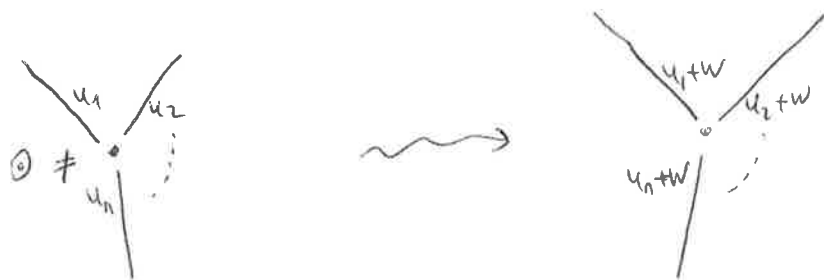
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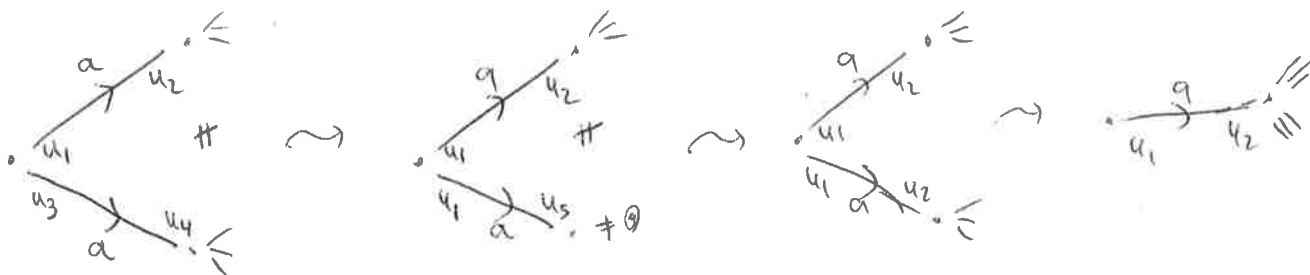
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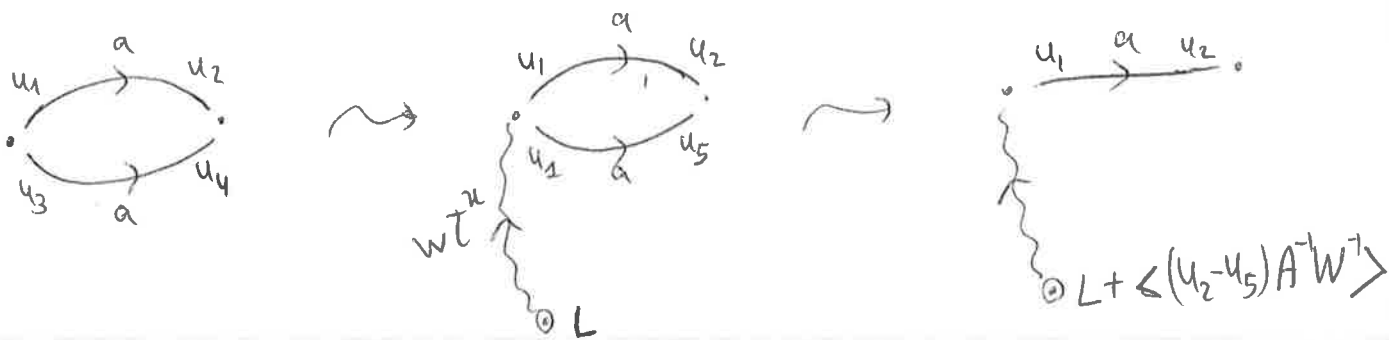
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open folding



closed folding



$$\begin{aligned}
 & w t^u t^{-u_1} a t^{u_2} t^{-u_5} a^{-1} t^{u_4} t^{-u} W^{-1} = \\
 & = \cancel{w a a^{-1} W^{-1}} \cdot t^{(u_2-u_1)W^{-1} + (u_2-u_5)A^{-1}W^{-1} + (u_4-u_1)W^{-1}} \\
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Lemma

With repeated use of the above operations, any vectored X -automata \mathcal{A} can be converted into a vectored Stallings X -automata \mathcal{A}' .

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For $H \leq \mathbb{G}_{\mathbf{A}}$, the result of folding $\mathcal{F}I(H)$, say $\Gamma(H)$, is uniquely determined by the subgroup H , modulo the choice of the maximal tree, and with all vectors around understood 'modulo L '.

Theorem (Delgado–V., 2016)

The following map is a bijection:

$$\begin{array}{ccc}
 \{\text{Subgroups of } \mathbb{G}_{\mathbf{A}}\} & \longleftrightarrow & \{\text{Vectored Stallings } X\text{-automata}\} \\
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Membership

Corollary (Delgado–V., 2016)

Membership is solvable in free-abelian-by-free groups.

Proof. Given $wt^u \in \mathbb{G}_A$ and $H = \langle w_1 t^{u_1}, \dots, w_r t^{u_r} \rangle \leq \mathbb{G}_A$, do the following:

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The Subgroup Intersection Problem

(Caution!)

The groups \mathbb{G}_A are **NOT** Howson, in general.

Definition (The Subgroup Intersection Problem for G)

Input: $g_1, \dots, g_r, g'_1, \dots, g'_s \in G$
Decide: $\langle g_1, \dots, g_r \rangle \cap \langle g'_1, \dots, g'_s \rangle$ is f.g. and, if so, compute generators.

Theorem (Delgado–V., 2017)

The Subgroup Intersection Problem is solvable in $\mathbb{F}_n \times \mathbb{Z}^m$.

Question

What about \mathbb{G}_A ?

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Given $H, K \leq_{fg} \mathbb{G}_A$, we know that

$$H \cap K \text{ is f.g.} \Leftrightarrow (H \cap K)\pi \text{ is f.g.}$$

and $H\pi \cap K\pi$ is f.g. (because \mathbb{F}_n is Howson) ... **BUT**

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with possibly **strict** inequality.

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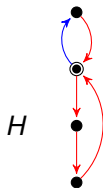
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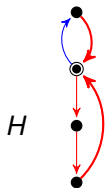
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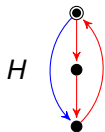
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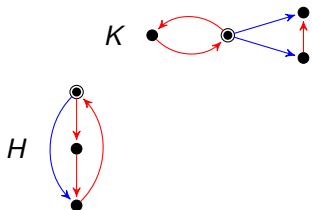
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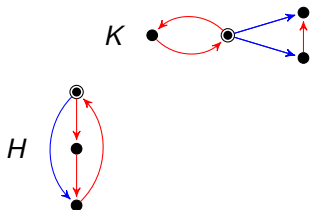
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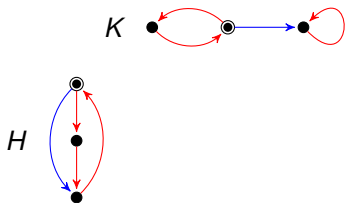
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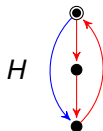
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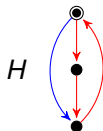
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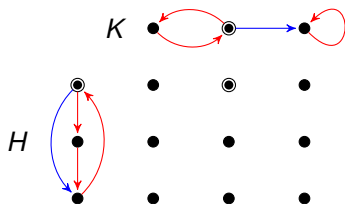
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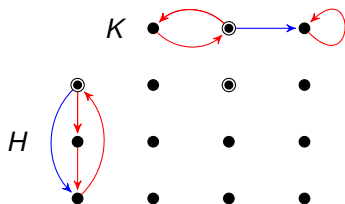
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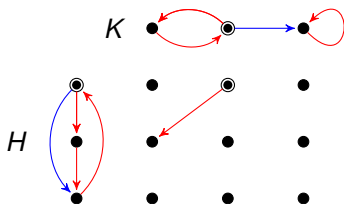
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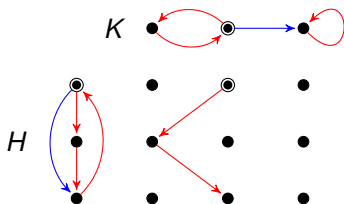
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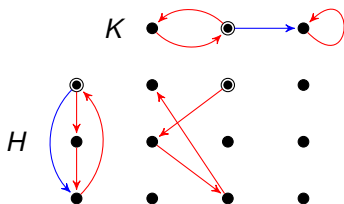
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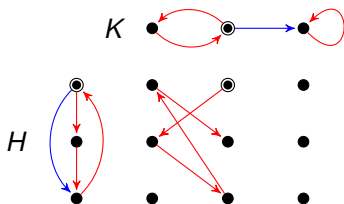
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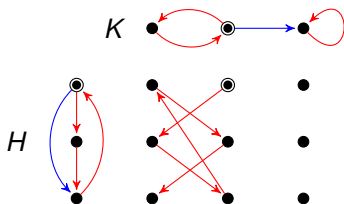
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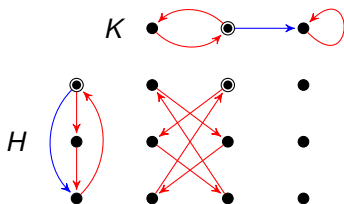
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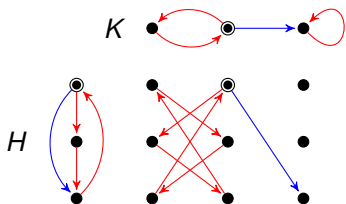
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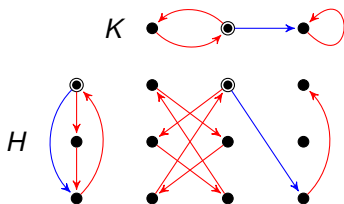
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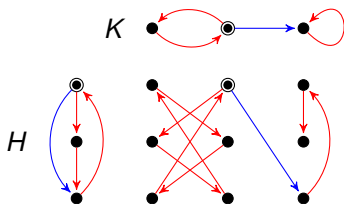
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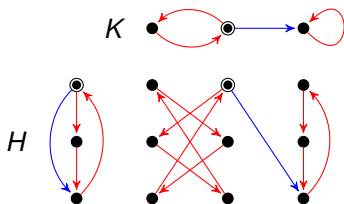
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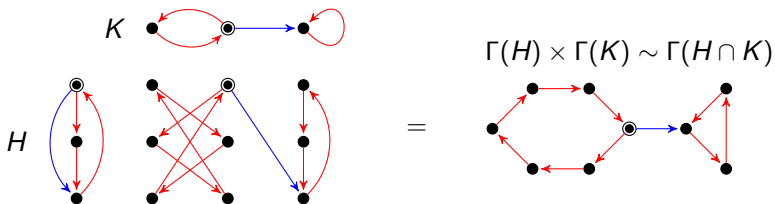
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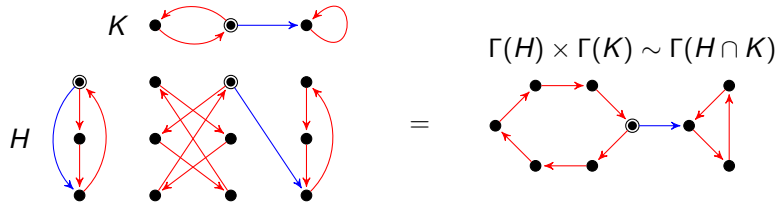
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Therefore, $H \cap K = \langle x^6, yx^3y^{-1} \rangle$.

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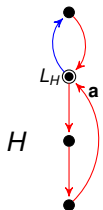
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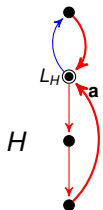
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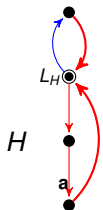
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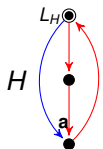
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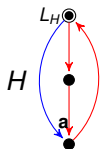
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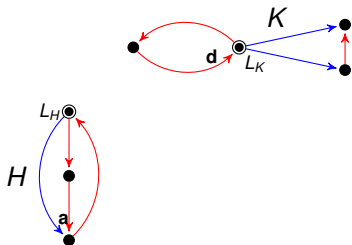
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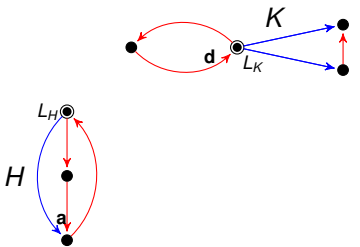
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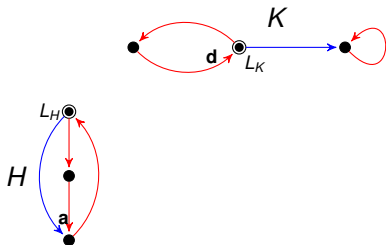
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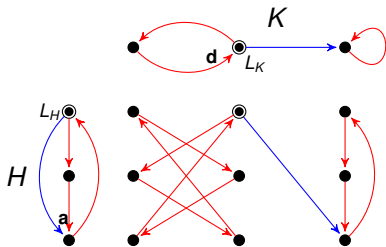
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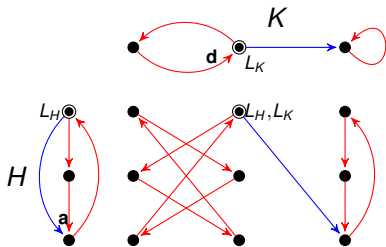
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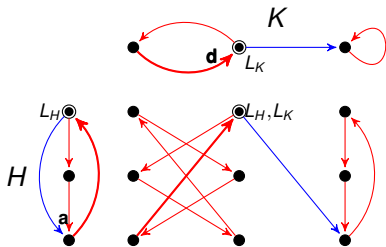
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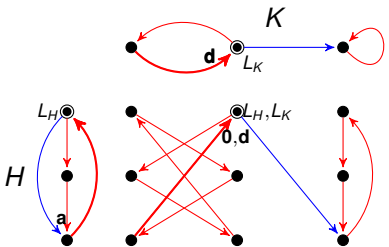
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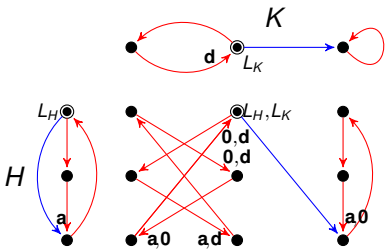
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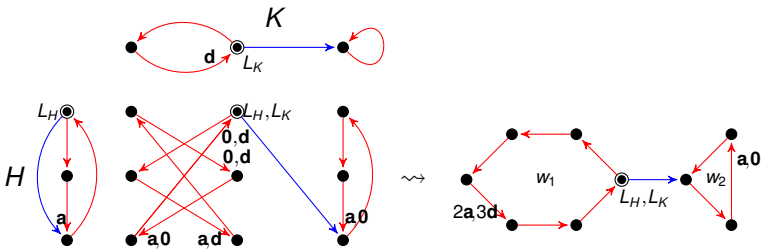
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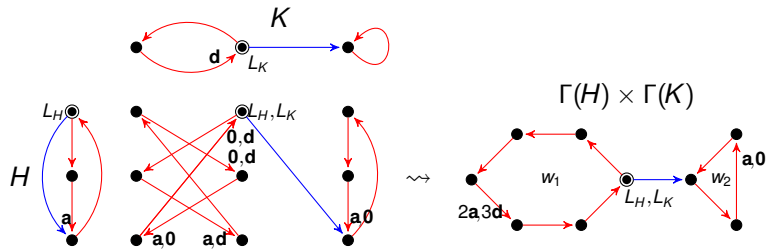
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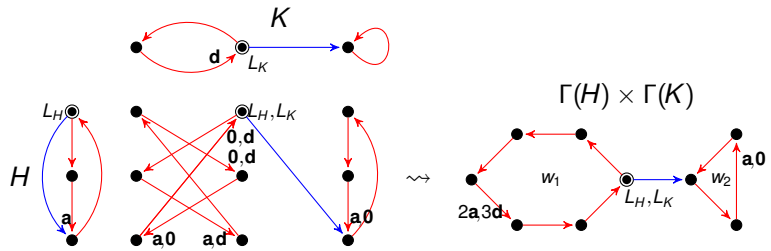
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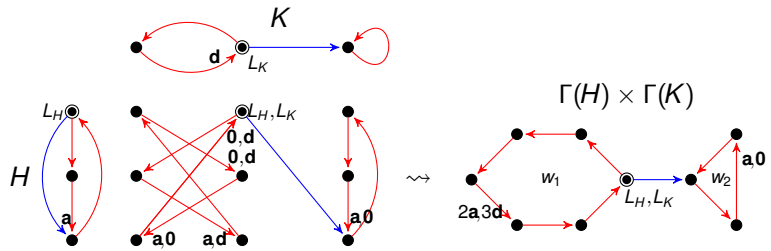
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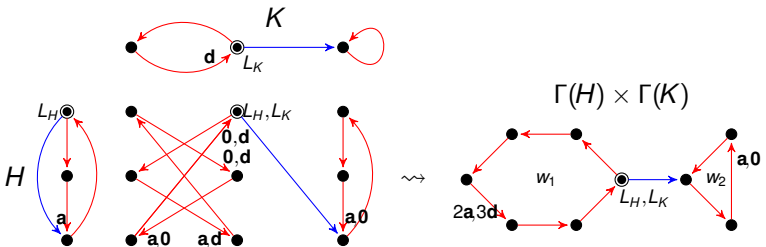
Claim:

$$H \cap K = \langle ut^a \mid ut^a \text{ is componentwise-readable in } \Gamma(H) \times \Gamma(K) \rangle$$

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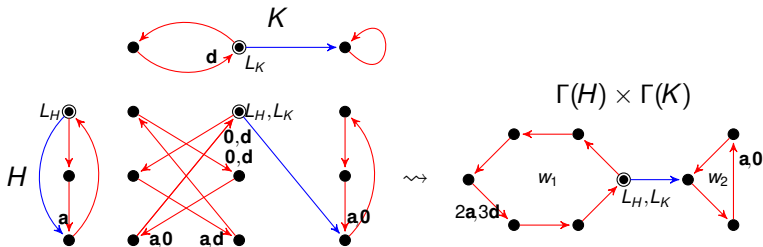
$$H \cap K = \langle ut^a \mid ut^a \text{ is componentwise-readable in } \Gamma(H) \times \Gamma(K) \rangle$$

$$(H \cap K)_\pi = \langle w \in \mathbb{F}_{\{w_1, w_2\}} \mid w(w_1 t^{2a}, w_2 t^a) t^{L_H} \cap w(w_1 t^{3d}, w_2 t^0) t^{L_K} \neq \emptyset \rangle$$

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Intersections in $\mathbb{F}_n \times \mathbb{Z}^m$

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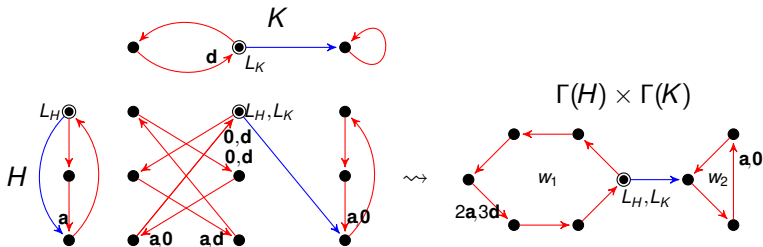
Claim:

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$$\begin{aligned} (H \cap K)\pi &= \langle w \in \mathbb{F}_{\{w_1, w_2\}} \mid w(w_1 t^{2a}, w_2 t^a) t^{L_H} \cap w(w_1 t^{3d}, w_2 t^0) t^{L_K} \neq \emptyset \rangle \\ &= \langle w \in \mathbb{F}_{\{w_1, w_2\}} \mid w^{ab} \begin{bmatrix} 2a-3d \\ a-0 \end{bmatrix} \in L_H + L_K \rangle \end{aligned}$$

Intersections in $\mathbb{F}_n \times \mathbb{Z}^m$

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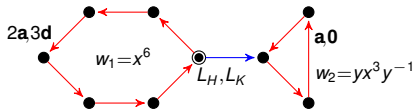
$$H \cap K = \langle ut^a \mid ut^a \text{ is componentwise-readable in } \Gamma(H) \times \Gamma(K) \rangle$$

$$(H \cap K)\pi = \left\langle w \in \mathbb{F}_{\{w_1, w_2\}} \mid w(w_1 t^{2a}, w_2 t^a) t^{L_H} \cap w(w_1 t^{3d}, w_2 t^0) t^{L_K} \neq \emptyset \right\rangle$$

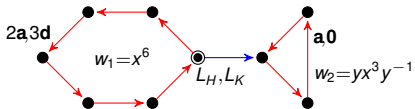
$$= \left\langle w \in \mathbb{F}_{\{w_1, w_2\}} \mid w^{ab} \begin{bmatrix} 2a-3d \\ a-0 \end{bmatrix} \in L_H + L_K \right\rangle$$

$$= (L_H + L_K) \mathbf{B}^{-1} \rho^{-1}, \text{ where } \mathbf{B} = \begin{bmatrix} 2a-3d \\ a-0 \end{bmatrix} \text{ and } \rho = ab.$$

Intersections in $\mathbb{F}_n \times \mathbb{Z}^m$



Intersections in $\mathbb{F}_n \times \mathbb{Z}^m$

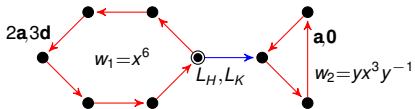


$$\mathbf{B} = \begin{bmatrix} 2a-3d \\ a-0 \end{bmatrix}$$

$$M = (L_H + L_K)\mathbf{B}^{-1}$$

We have that $(H \cap K)\pi = (L_H + L_K)\mathbf{B}^{-1}\rho^{-1} = M\rho^{-1}$, i.e.,

Intersections in $\mathbb{F}_n \times \mathbb{Z}^m$



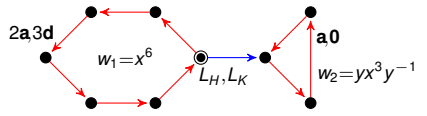
$$\mathbf{B} = \begin{bmatrix} 2\mathbf{a}-3\mathbf{d} \\ \mathbf{a}-\mathbf{0} \end{bmatrix}$$

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$$\begin{array}{ccccccc} \mathbb{F}_2 \supseteq & H\pi \cap K\pi & = & \mathbb{F}_{\{w_1, w_2\}} & \xrightarrow{\rho} & \mathbb{Z}^2 & \xrightarrow{\mathbf{B}} & \mathbb{Z}^2 \\ & \downarrow & & \downarrow & & \downarrow & & \downarrow \\ & 1 \neq (H \cap K)\pi & = & M\rho^{-1} & \longleftarrow & M & \longleftarrow & L_H + L_K \end{array}$$

Intersections in $\mathbb{F}_n \times \mathbb{Z}^m$



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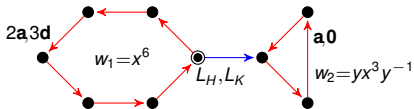
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We have that $(H \cap K)\pi = (L_H + L_K)\mathbf{B}^{-1}\rho^{-1} = M\rho^{-1}$, i.e.,

$$\begin{array}{ccccccc} \mathbb{F}_2 \geq & H\pi \cap K\pi & = & \mathbb{F}\{w_1, w_2\} & \xrightarrow{\rho} & \mathbb{Z}^2 & \xrightarrow{\mathbf{B}} & \mathbb{Z}^2 \\ & \nabla & & \nabla & & \nabla & & \nabla \\ & 1 \neq (H \cap K)\pi & = & M\rho^{-1} & \longleftarrow & M & \longleftarrow & L_H + L_K \end{array}$$

Then, $\Gamma((H \cap K)\pi, \{w_1, w_2\}) = \Gamma(M\rho^{-1}, \{w_1, w_2\})$

Intersections in $\mathbb{F}_n \times \mathbb{Z}^m$



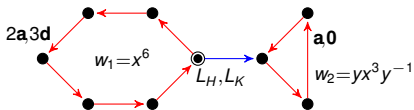
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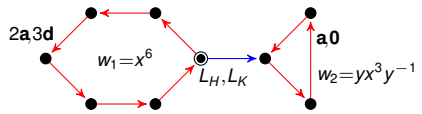
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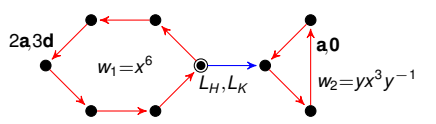
$\nabla \qquad \qquad \nabla \qquad \qquad \nabla \qquad \qquad \nabla$

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$$\begin{aligned}
 &= \text{Sch}(M\rho^{-1}, \{w_1, w_2\}) \\
 &= \text{Cay}(\mathbb{F}_{\{w_1, w_2\}}/M\rho^{-1}, \{[w_1], [w_2]\}) \\
 &= \text{Cay}(\mathbb{Z}^2/M, \{\mathbf{e}_1, \mathbf{e}_2\})
 \end{aligned}$$

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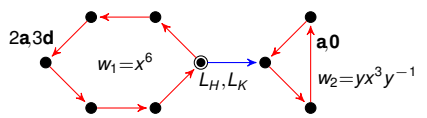
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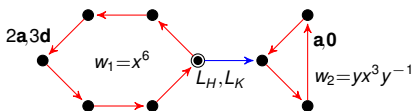
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$$= \text{Sch}(M\rho^{-1}, \{w_1, w_2\})$$

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$$= \text{Cay}(\mathbb{Z}^2/\text{row}(\mathbf{M}), \{\mathbf{e}_1, \mathbf{e}_2\})$$

Intersections in $\mathbb{F}_n \times \mathbb{Z}^m$



$$\mathbf{B} = \begin{bmatrix} 2a-3d \\ a-0 \end{bmatrix}$$

$$\langle \mathbf{M} \rangle = M = (L_H + L_K) \mathbf{B}^{-1}$$

$$\mathbf{PMQ} = \mathbf{D} = \text{diag}(\delta_1, \delta_2)$$

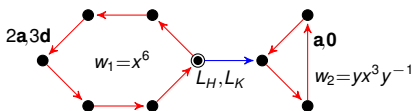
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$$\quad \quad \quad \nabla \quad \quad \quad \nabla \quad \quad \quad \nabla \quad \quad \quad \nabla$$

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Intersections in $\mathbb{F}_n \times \mathbb{Z}^m$ 

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Intersection example: case 1

$$H = \langle t^{L_H}, x^3 t^a, yx \rangle, K = \langle t^{L_K}, x^2 t^d, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

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$$H = \langle t^{L_H}, x^3 t^{\mathbf{a}}, yx \rangle, K = \langle t^{L_K}, x^2 t^{\mathbf{d}}, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

Case 1: $\mathbf{a} = (1, 0)$, $\mathbf{d} = (0, 1)$, $L_H = \langle (0, 6) \rangle$, $L_K = \langle (3, -3) \rangle$.

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$$H = \langle t^{L_H}, x^3 t^{\mathbf{a}}, yx \rangle, K = \langle t^{L_K}, x^2 t^{\mathbf{d}}, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

Case 1: $\mathbf{a} = (1, 0)$, $\mathbf{d} = (0, 1)$, $L_H = \langle (0, 6) \rangle$, $L_K = \langle (3, -3) \rangle$.

Then, $\mathbf{B} = \begin{bmatrix} 2 & -3 \\ 1 & 0 \end{bmatrix}$, $\mathbf{M} = \begin{bmatrix} -2 & 4 \\ 1 & 1 \end{bmatrix}$, $\mathbf{Q} = \begin{bmatrix} 1 & -1 \\ 0 & 1 \end{bmatrix}$, $\mathbf{D} = \begin{bmatrix} 1 & 0 \\ 0 & 6 \end{bmatrix}$.

Intersection example: case 1

$$H = \langle t^{L_H}, x^3 t^{\mathbf{a}}, yx \rangle, K = \langle t^{L_K}, x^2 t^{\mathbf{d}}, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

Case 1: $\mathbf{a} = (1, 0)$, $\mathbf{d} = (0, 1)$, $L_H = \langle (0, 6) \rangle$, $L_K = \langle (3, -3) \rangle$.

Then, $\mathbf{B} = \begin{bmatrix} 2 & -3 \\ 1 & 0 \end{bmatrix}$, $\mathbf{M} = \begin{bmatrix} -2 & 4 \\ 1 & 1 \end{bmatrix}$, $\mathbf{Q} = \begin{bmatrix} 1 & -1 \\ 0 & 1 \end{bmatrix}$, $\mathbf{D} = \begin{bmatrix} 1 & 0 \\ 0 & 6 \end{bmatrix}$. Hence:

$$\Gamma((H \cap K)\pi, \{w_1, w_2\}) = \text{Cay}(\mathbb{Z}/1\mathbb{Z} \oplus \mathbb{Z}/6\mathbb{Z}, \{(1, -1), (0, 1)\})$$

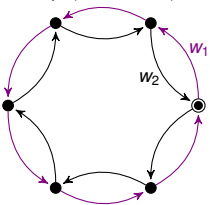
Intersection example: case 1

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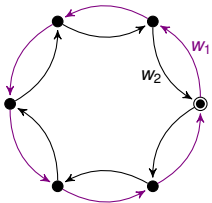


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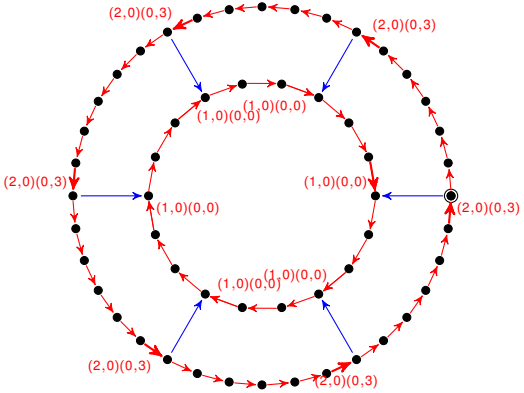
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Intersection example: case 1

$$H = \langle t^{L_H}, x^3 t^a, yx \rangle, K = \langle t^{L_K}, x^2 t^d, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

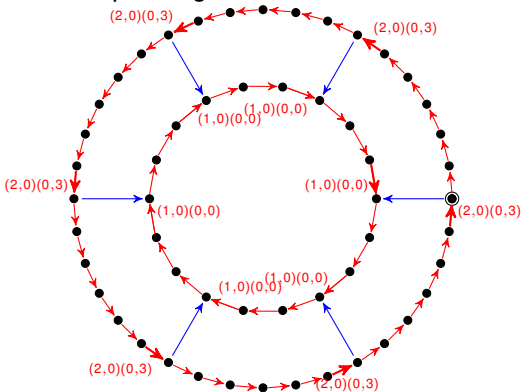
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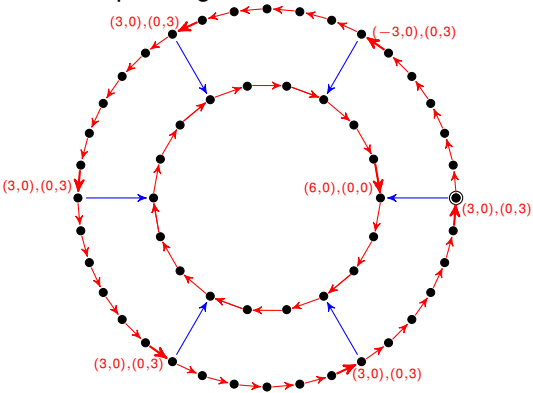
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Intersection example: case 1

$$H = \langle t^{L_H}, x^3 t^a, yx \rangle, K = \langle t^{L_K}, x^2 t^d, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

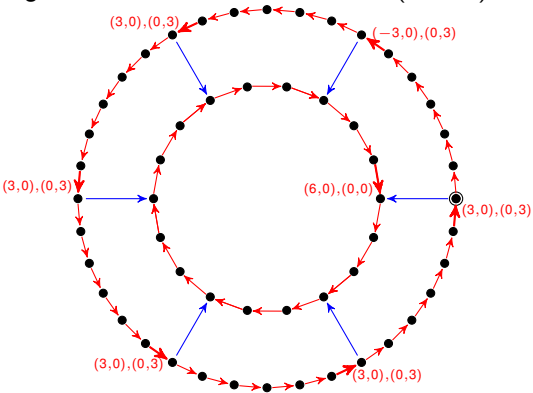
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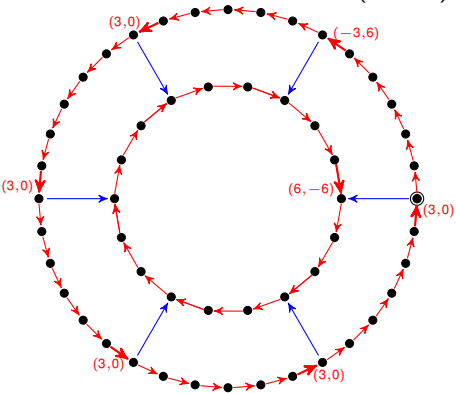
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 after equalizing the abelian labels we obtain $\Gamma(H \cap K)$:



Intersection example: case 1

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Intersection example: case 1

Therefore, $H = \langle t^{(0,6)}, x^3 t^{(1,0)}, yx \rangle$ $K = \langle t^{(3,-3)}, x^2 t^{(0,1)}, yxy^{-1} \rangle$ and

$$H \cap K = \left\langle \begin{array}{l} x^6 y x^3 y^{-1} t^{(3,0)}, \\ x^{12} y x^3 y^{-1} x^{-6} t^{(3,0)}, \\ x^{18} y x^3 y^{-1} x^{-12} t^{(3,0)}, \\ x^{24} y x^3 y^{-1} x^{-18} t^{(3,0)} \end{array}, \begin{array}{l} x^{30} y x^3 y^{-1} x^{-24} t^{(3,0)}, \\ x^{36} t^{(12,6)}, \\ y x^{18} y^{-1} t^{(6,-6)}, \end{array} \right\rangle$$

Note, for example, that

$$H \cap K \ni x^{36} t^{(12,6)} = \begin{cases} = (x^3 t^{(1,0)})^{12} t^{(0,6)} \in H, \\ = (x^2 t^{(0,1)})^{18} (t^{(3,-3)})^4 \in K. \end{cases}$$

And that $x^6 \in H\pi \cap K\pi$ but $x^6 \notin (H \cap K)\pi$ since

$$x^6 t^{(2,0) + \lambda(0,6)} \in H$$

$$x^6 t^{(0,3) + \mu(3,-3)} \in K$$

but $\left((2,0) + \langle (0,6) \rangle \right) \cap \left((0,3) + \langle (3,-3) \rangle \right) = \emptyset$.

Intersection example: case 2

$$H = \langle t^{L_H}, x^3 t^a, yx \rangle, K = \langle t^{L_K}, x^2 t^d, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

Intersection example: case 2

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Case 2: $\mathbf{a} = (3, 3)$, $\mathbf{d} = (2, 2)$, $L_H = \langle (1, 2) \rangle$, $L_K = \langle (0, 0) \rangle$.

Intersection example: case 2

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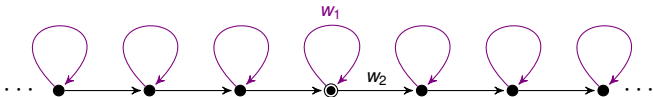
Then, $\delta_1 = 1$, $\delta_2 = 0$ and so, $\Gamma(H \cap K) = \text{Cay}(\mathbb{Z}, \{0, 1\})$.

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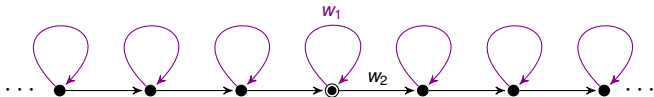


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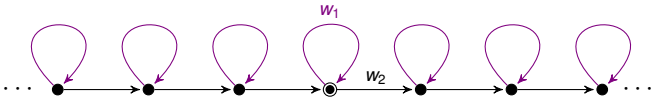
After replacing, folding, normalizing, and equalizing, we obtain $\Gamma(H \cap K)$:

Intersection example: case 2

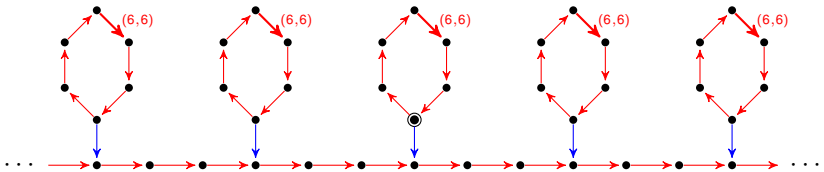
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Then, $\delta_1 = 1, \delta_2 = 0$ and so, $\Gamma(H \cap K) = \text{Cay}(\mathbb{Z}, \{0, 1\})$.



After replacing, folding, normalizing, and equalizing, we obtain $\Gamma(H \cap K)$:



Intersection example: case 3

$$H = \langle t^{L_H}, x^3 t^a, yx \rangle, K = \langle t^{L_K}, x^2 t^d, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

Intersection example: case 3

$$H = \langle t^{L_H}, x^3 t^{\mathbf{a}}, yx \rangle, K = \langle t^{L_K}, x^2 t^{\mathbf{d}}, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

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Intersection example: case 3

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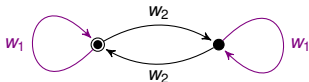
Then, $\delta_1 = 1$, $\delta_2 = 2$ and so, $\Gamma(H \cap K) = \text{Cay}(\mathbb{Z}/2\mathbb{Z}, \{0, 1\})$.

Intersection example: case 3

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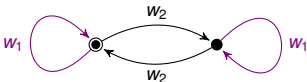


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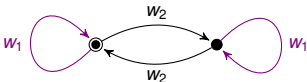


After replacing, folding, normalizing, and equalizing, we obtain $\Gamma(H \cap K)$:

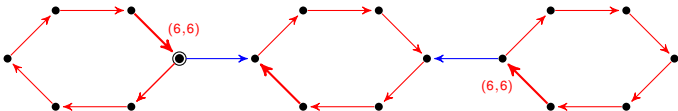
Intersection example: case 3

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After replacing, folding, normalizing, and equalizing, we obtain $\Gamma(H \cap K)$:



Intersection example: case 4

$$H = \langle t^{L_H}, x^3 t^a, yx \rangle, K = \langle t^{L_K}, x^2 t^d, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

Intersection example: case 4

$$H = \langle t^{L_H}, x^3 t^{\mathbf{a}}, yx \rangle, K = \langle t^{L_K}, x^2 t^{\mathbf{d}}, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

Case 4: $\mathbf{a} = (6, 6)$, $\mathbf{d} = (4, 4) \in \mathbb{Z}^2$, $L_H = \langle (6p, 6p) \rangle$, $L_K = \langle (0, 0) \rangle$,
for some $p \in \mathbb{Z}$, $p \neq 0$.

Intersection example: case 4

$$H = \langle t^{L_H}, x^3 t^{\mathbf{a}}, yx \rangle, K = \langle t^{L_K}, x^2 t^{\mathbf{d}}, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

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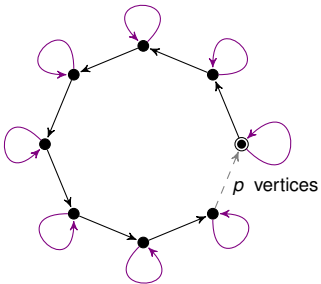
Then, $\delta_1 = 1$, $\delta_2 = p$ and so, $\Gamma(H \cap K) = \text{Cay}(\mathbb{Z}/p\mathbb{Z}, \{0, 1\})$

Intersection example: case 4

$$H = \langle t^{L_H}, x^3 t^{\mathbf{a}}, yx \rangle, K = \langle t^{L_K}, x^2 t^{\mathbf{d}}, yxy^{-1} \rangle \leq \mathbb{F}_2 \times \mathbb{Z}^2$$

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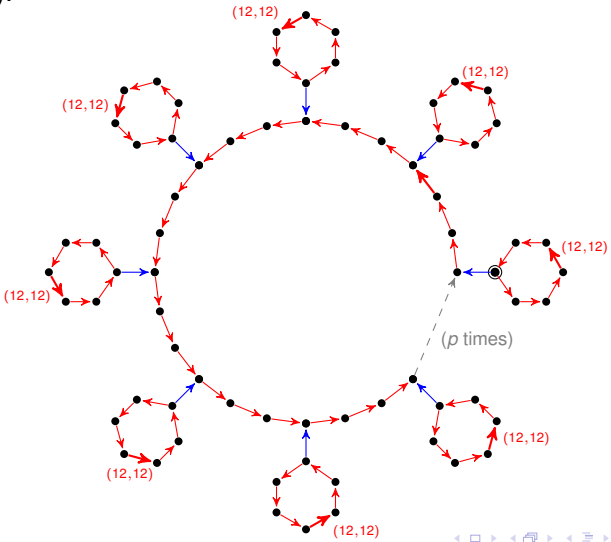


Intersection example: case 4

After replacing, folding, normalizing, and equalizing, we obtain $\Gamma(H \cap K)$:

Intersection example: case 4

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Intersection example: case 4

Therefore, $H_p = \langle t^{(6p,6p)}, x^3 t^{(6,6)}, yx \rangle$, $K = \langle x^2 t^{(4,4)}, yxy^{-1} \rangle$ and

$$H \cap K = \left\langle \begin{array}{c} (yx^{-3}y^{-1})^{-1} x^6 (yx^{-3}y^{-1}) t^{(12,12)} \\ (yx^{-3}y^{-1})^{-2} x^6 (yx^{-3}y^{-1})^2 t^{(12,12)} \\ \vdots \\ (yx^{-3}y^{-1})^{-(p-1)} x^6 (yx^{-3}y^{-1})^{p-1} t^{(12,12)} \\ yx^{3p}y^{-1} \end{array} \right\rangle$$

Note that $r(H_p) = 3$, $r(K) = 2$, but

$$r(H_p \cap K) = p + 1.$$

So, **no** Hanna Neumann type inequality $\tilde{r}(H \cap K) \leq C \cdot \tilde{r}(H) \cdot \tilde{r}(K)$ is possible in these groups.

THANKS